

## PATH-DECOMPOSITIONS

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A path-decomposition of a graph is a partition of its edges into subgraphs each of which is a path or a union of paths (a linear forest). We survey known results when the graph and the linear forest are of prescribed types, and when the decomposition satisfies further additional properties.

### 1. Introduction.

The basic question concerning graph decomposition is: *When can the edges of a given graph be partitioned into a specified collection of subgraphs?* Specifically, let  $G$  be a graph and  $\mathcal{F} = \{F^1, F^2, \dots, F^r\}$  be a family of graphs. An  $\mathcal{F}$ -decomposition of  $G$  (also called a  $(G, \mathcal{F})$ -design) is a partition of the edge-set of  $G$  into the  $r$  subgraphs  $F^1, F^2, \dots, F^r$ . In the case when all subgraphs in  $\mathcal{F}$  are isomorphic to the graph  $F$  we will refer to the design as  $(G, F)$ -design. To date the most far-reaching result is that of Wilson [71] who proved that there is a  $(K_n, F)$ -design for all sufficiently large  $n$  satisfying  $|V(F)| \leq n, n(n-1) \equiv 0 \pmod{2|E(F)|}$  and  $n-1 \equiv 0 \pmod{d}$ , where  $d$  is the greatest common divisor of the degrees of the vertices in  $F$ . (Note that except for the "size" of  $n$  these are precisely the easily determined necessary conditions for the existence of such a design.) Several surveys of  $(G, \mathcal{F})$ -designs have been written and the reader is referred to Akiyama and Kano [1], Bermond and Sotteau [10], Chung and Graham [18], Harary and Robinson [26], Rodger [58] and the outstanding book of Bosák [13].

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The particular question we will consider in this paper is: For what  $G$  and  $\mathcal{F}$ , where the subgraphs in  $\mathcal{F}$  are linear forests, does a  $(G, \mathcal{F})$ -design exist? (A *linear forest* is a graph each component of which is a path.) Throughout we will always assume the graph being considered is simple unless it is stated otherwise.

What follows is a list of notation that will be used:

$G^\lambda$	The graph $G$ in which every edge has multiplicity $\lambda$ .
$DG$	The graph $G$ with each edge replaced by two oppositely directed arcs.
$mG$	$m$ edge-disjoint copies of $G$ .
$E(G)$	The set of edges in the graph $G$ .
$V(G)$	The set of all vertices in $G$ .
$\chi'(G)$	The edge-chromatic number of $G$ .
$\Delta(G)$	The maximum degree in $G$ .
$\delta(G)$	The minimum degree in $G$ .
$G \cup H$	The edge-disjoint union of the graphs $G$ and $H$ .
$K_n$	The complete graph on $n$ vertices.
$P_k$	The path with $k$ vertices and $k - 1$ edges.
$\vec{P}_k$	The directed path $P_k$ .
$C_k$	The cycle with $k$ vertices.
$K_{(m,n)}$	The complete multipartite graph with $m$ parts each of size $n$ .
$K_{a_1, a_2, \dots, a_m}$	The complete multipartite graph with $m$ parts of sizes $a_1, a_2, \dots, a_m$ .
$Z_m$	The cyclic group of order $m$ .

## 2. $(G, F)$ -designs, where $F$ is a linear forest.

A  $(G, F)$ -design when  $F$  is a path is usually referred to as a path-decomposition. Certainly the most well known path-decomposition is the  $(K_{2n}, P_{2n})$ -design which is easily derived from the  $(K_{2n}, C_{2n})$ -design (given by Walecki [51, page 162]) on deleting a vertex.

**Theorem 2.1.** *There exists a  $(K_{2n}, P_{2n})$ -design for all values of  $n$ .*

*Proof.* With the vertices of  $K_{2n}$  labelled by  $Z_{2n}$ , the paths are  $F^i = [i, i + 1, i + 2n - 1, i + 2, \dots, i + n + 2, i + n - 1, i + n + 1, i + n]$ ,  $1 \leq i \leq n$ .

Many years later Hell and Rosa [30] defined a *balanced graph design* to be a  $(K_n^\lambda, F)$ -design with the additional property that each vertex of  $K_n^\lambda$  lies in the same number of copies of  $F$  (hence the term balanced). They focussed their

attention on the case when  $F \cong P_k$ . It is not difficult to show the existence of all  $(K_n^\lambda, P_n)$ -designs (which are necessarily balanced).

**Lemma 2.2.** *There exists a  $(K_{2n+1}^2, P_{2n+1})$ -design for all values of  $n$ .*

*Proof.* With the vertices of  $K_{2n+1}^2$  labelled by  $Z_{2n+1}$ , the paths are  $F^i = [i, i + 1, i + 2n, i + 2, i + 2n - 1, \dots, i + n - 1, i + n + 2, i + n, i + n + 1]$ ,  $1 \leq i \leq 2n + 1$ .

**Corollary 2.3.** *There exists a balanced  $(K_n^\lambda, P_n)$ -design if and only if  $\lambda n$  is even.*

Partial results on the existence of balanced path-decompositions were obtained by Hell and Rosa [30], Hung and Mendelsohn [40], and Lawless [46, 47]. Soon after complete solutions were independently presented by Huang [36] and Hung and Mendelsohn [41].

**Theorem 2.4.** (Huang; Hung and Mendelsohn) *There exists a balanced  $(K_n^\lambda, P_k)$ -design if and only if*

$$\lambda n (n - 1) \equiv 0 \pmod{2(k - 1)}$$

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Huang [37] also gave necessary and sufficient conditions for the existence of balanced  $(K_n^\lambda, P_2 \cup P_3)$ -designs, where  $P_2 \cup P_3$  is a linear forest.

Others were meanwhile considering the more general question of the existence of  $(K_n^\lambda, F)$ -designs, particularly for small graphs  $F$ . When  $F$  is the path  $P_k$  it is easy to see that the only necessary condition for the existence of the design is  $\lambda n (n - 1) \equiv 0 \pmod{2(k - 1)}$ . The existence of these designs is trivially observed when  $k = 2$ . We will prove the result for the case  $k = 3$  and, although the proof is quite simple, it serves to illustrate the recursive technique that is frequently used.

**Theorem 2.5.** *There exists a  $(K_n^\lambda, P_3)$ -design if and only if  $\lambda n (n - 1) \equiv 0 \pmod{4}$ .*

*Proof.* Only two cases need to be considered:

(a)  $\lambda = 1$  and  $n \equiv 0, 1 \pmod{4}$  and

(b)  $\lambda = 2$  and  $n \equiv 2, 3 \pmod{4}$ .

In all constructions  $V(K_{2m+1}^\lambda) = Z_{2m+1}$  and  $V(K_{2m}^\lambda) = Z_{2m-1} \cup \{\infty\}$ . We begin with the small designs.

$(K_4, P_3)$ -design:  $F^i = [\infty, i, i + 1]$ ,  $1 \leq i \leq 3$

$(K_5, P_3)$ -design:  $F^i = [i, i + 1, i + 3]$ ,  $1 \leq i \leq 5$

$(K_6^2, P_3)$ -design:  $F^i = [i, \infty, i + 1]$ ,  $F^{i+5} = [i, i + 1, i + 2]$ ,  $F^{i+10} = [i, i + 2, i + 4]$ ,  $1 \leq i \leq 5$

$(K_7^2, P_3)$ -design:  $F^i = [i, i + 1, i + 2]$ ,  $F^{i+7} = [i, i + 2, i + 4]$ ,  $F^{i+14} = [i, i + 3, i + 6]$ ,  $1 \leq i \leq 7$ .

Also note that  $C_4 \cong K_{2,2}$  is the union of two paths of length two.

- (a) If  $n = 4t$ , write  $K_{4t} = tK_4 \cup \binom{t}{2}K_{4,4}$ . Since each subgraph has a decomposition into paths of length two so too does  $K_{4t}$ . In the case  $n = 4t + 1$  add a new vertex,  $\infty$ , and replace each of the  $t(K_4, P_3)$ -designs with a  $(K_5, P_3)$ -design on the original four vertices and the vertex  $\infty$ .
- (b) If  $n = 4t + 2$ , write  $K_{4t+2}^2 = (t - 1)K_4^2 \cup K_6^2 \cup \binom{t-1}{2}K_{4,4}^2 \cup (t - 1)K_{4,6}^2$ . Again, as each subgraph has a decomposition into paths of length two so too does  $K_{4t+2}^2$ . In the case  $n = 4t + 3$  add a new vertex,  $\infty$ , as was done in part (a), noting that a  $(K_7^2, P_3)$ -design is also required.

Bermond and Schönheim [9] gave constructions for  $(K_n, P_4)$ -designs ( $n \equiv 0, 1 \pmod{3}$ ) and Huang and Rosa [39] constructed  $(K_n, P_k)$ -designs,  $2 \leq k \leq 9$ , for all possible values of  $n$ . Bermond, Huang, Rosa and Sotteau [8] also constructed all  $(K_n, P_5)$ -designs and all  $(K_n, P_2 \cup P_3)$ -designs. To construct  $(K_n, 2P_3)$ -designs is a straightforward exercise. In fact using the recursive method illustrated in Theorem 2.5, it is not difficult (although it does take a little time) to construct all  $(K_n^\lambda, F)$ -designs, where  $F$  is a linear forest on at most six vertices.

Finally in 1963, Tarsi [65] completely settled the question of the existence of  $(K_n^\lambda, P_k)$ -designs.

**Theorem 2.6.** (Tarsi) *A  $(K_n^\lambda, P_k)$ -design exists if and only if  $\lambda n(n - 1) \equiv 0 \pmod{2(k - 1)}$  and  $n \geq k$ .*

His method is quite different from either of the two we have so far described. For  $k \leq n - 2$ , when  $\lambda(n - 1)$  is even,  $K_n^\lambda$  is Eulerian. Tarsi produced an Eulerian walk in which any two occurrences of a particular vertex are separated by at least  $k - 1$  distinct vertices all different from it. This path can then be "broken" into copies of  $P_k$ . When  $\lambda(n - 1)$  is odd a collection of paths of length  $k$  is first chosen so that no edge occurs in them more than  $\lambda$  times and each vertex of  $K_n^\lambda$  is an endpoint of exactly one of them. When these paths are deleted from the graph what remains is Eulerian and again an appropriate Eulerian walk is described. The case  $k = n$  and  $k = n - 1$  are treated separately (but note Corollary 2.3 provides a solution in the case  $k = n$ ).

Let us now return to the paper of Hell and Rosa. The main goal of that paper

was to find all  $(K_n^\lambda, \frac{n}{3}P_3)$ -designs ( $n \equiv 0 \pmod{3}$ ), where  $\frac{n}{3}P_3$  is a linear forest with each component a path of length 2. The question of the existence of such designs arose from a recreational mathematics problem posed by Dudeney [19, Problem #272] in which he asks the reader to arrange nine schoolboys so that on each of six days they walk in three rows of three and on no occasion does any boy walk twice beside any other. (The naming of the problem as the "handcuffed prisoners problem" is a result of an alternate formulation also given by Dudeney [20].)

To simplify notation, we will refer to a  $(K_n^\lambda, (n/k)P_k)$ -design (where necessarily  $n \equiv 0 \pmod{k}$ ), as a *resolvable*  $(K_n^\lambda, P_k)$ -design; meaning that we have a  $(K_n^\lambda, P_k)$ -design with the property that the paths can be partitioned into vertex-disjoint spanning subgraphs of  $K_n^\lambda$ . Each spanning subgraph is referred to as a *factor*. (Such designs are necessarily balanced).

Hell and Rosa [30] and Nakamura, Kimura and Suganuma [53] obtained partial results on the existence of these designs for general  $k$ . The second authors considered in particular the case  $k = 3$  and in an unpublished manuscript (in Japanese) they describe an investigation of the general case. Huang and Rosa [38] made a thorough study of the resolvable  $(K_9, P_3)$ -designs showing that there are precisely 334 of them which are not isomorphic. Fortunately, this calculation was repeated by Ollerenshaw and Bondi [55] who determined that in fact there are 332 non-isomorphic designs. The case  $k = 3$  was finally resolved by Horton [33].

**Theorem 2.7.** (Horton) *There exists a resolvable  $(K_n^\lambda, P_3)$ -design if and only if  $n \equiv 0 \pmod{3}$  and  $\lambda(n-1) \equiv 0 \pmod{4}$ .*

In that paper Horton also showed that for  $n$  "sufficiently large" the necessary conditions for the existence of a  $(K_n, P_k)$ -design (obtained by simple counting) are also sufficient. The case  $k \geq 4$  was later resolved by Bermond, Heinrich and Yu [6].

**Theorem 2.8.** (Bermond, Heinrich and Yu) *There exists a resolvable  $(K_n^\lambda, P_k)$ -design ( $k \geq 4$ ) if and only if  $n \equiv 0 \pmod{k}$  and  $\lambda k(n-1) \equiv 0 \pmod{2(k-1)}$ .*

There are two somewhat surprising features of the proof. First, it is not particularly long and second, it does not contain the case  $k = 3$ . The proof relies on recursive ideas as described earlier and depends critically on the following lemma which appears to be extremely fruitful and which has been of use in other resolvable decomposition problems.

**Lemma 2.9.** *Let  $H$  be a multipartite graph with  $V(H) = \cup_{i=1}^k V_i$  so that for  $1 \leq i < j \leq k$ , the bipartite subgraph on vertex-set  $V_i \cup V_j$  with*

bipartition  $(V_i, V_j)$ , is  $t((i, j))$ -regular, where  $t$  is a mapping from the set  $\{(i, j) : 1 \leq i < j \leq k\}$  to the non-negative integers. Let  $G(H)$  be a graph with  $V(G(H)) = \{1, 2, \dots, k\}$  in which the edge  $ij$  has multiplicity  $t((i, j))$ . There is a resolvable  $(H, P_k)$ -design, if there is a  $(G(H), P_k)$ -design.

*Proof.* Suppose there is a  $(G(H), P_k)$ -design with paths  $P(1), P(2), \dots, P(m)$ , where  $m = \sum_{1 \leq i < j \leq k} \frac{t((i, j))}{(k-1)}$ . Each  $P(i)$  yields an appropriate factor of  $H$  as follows. To each edge  $pq$  of  $P(i)$  associate a 1-factor  $F_{pq}(i)$  from the bipartite subgraph with vertex-set  $V_p \cup V_q$ , so that  $\cup_{i=1}^m F_{pq}(i)$  is the bipartite subgraph on the vertex-set  $V_p \cup V_q$  (noting that  $F_{pq}(i)$  is empty if  $pq$  is not an edge of  $P(i)$ ). Clearly  $\cup_{pq \in P(i)} F_{pq}(i)$  is a factor of  $H$ , each component of which is a path of length  $k-1$ .

We next consider  $(G, F)$ -designs where  $G$  is other than a complete graph, and begin with  $F \cong P_3$ . Such designs were completely determined Kotzig [44] but we will present the proof of Caro and Schönheim [14]. (Note that the existence of a  $(G, P_3)$ -design is equivalent to the existence of a one-factor in the line graph of  $G$ , a question which had already been considered and resolved (for example, see [31]).)

**Theorem 2.10.** (Caro and Schönheim) *There is a  $(G, P_3)$ -design if and only if each component of  $G$  has an even number of edges.*

*Proof.* The necessity of the condition is obvious. The sufficiency is proved using induction on the number of edges in the graph and we need only consider the case when  $G$  is connected. If  $G$  has only two edges, then  $G \cong P_3$ . If  $G$  is a cycle of even length, there is a  $(G, P_3)$ -design. If  $G$  has no cut-edge and is not a cycle, choose an edge  $e$  of  $G$  with an adjacent edge  $f$  so that  $G - \{e, f\}$  is connected and apply the induction hypothesis to obtain a  $(G - \{e, f\}, P_3)$ -design.

If  $G$  has a cut-edge  $e$ , then  $G - e$  has two components  $H_1$  and  $H_2$ , where  $H_1$  has an even number of edges and  $H_2$  an odd number. If  $H_1$  is not a single vertex, then by induction we have an  $(H_1, P_3)$ -design and an  $(H_2 \cup e, P_3)$ -design. If  $H_1$  is a single vertex, let  $f$  be an edge of  $H_2$  which is adjacent with  $e$ . Then if  $f$  is not a cut-edge of  $H_2$  we have an  $(H_2 - f, P_3)$ -design and  $ef \cong P_3$ . If  $f$  is a cut-edge, then  $H_2 - f \cong H_3 \cup H_4$  and we either have an  $(H_3, P_3)$ -design and an  $(H_4, P_3)$ -design which together with the path  $ef$  give the decomposition, or an  $(H_3 \cup e, P_3)$ -design and a  $(H_4 \cup f, P_3)$ -design.

**Corollary 2.11.** *There is a  $(K_n^\lambda, P_3)$ -design if and only if  $\lambda n(n-1) \equiv 0 \pmod{4}$ .*

**Corollary 2.12.** *There is a  $(K_{n,m}^\lambda, P_3)$ -design if and only if  $\lambda mn$  is even and  $mn \geq 2$ .*

Truszczyński [67] considered the question of the existence of  $(K_{n,m}^\lambda, P_k)$ -designs.

**Theorem 2.13.** (Truszczyński) *If  $m \geq n$  and either  $\lambda$  is even or  $m$  and  $n$  are even, there is a  $(K_{m,n}^\lambda, P_k)$ -design if and only if  $\lambda mn \equiv 0 \pmod{k-1}$ ,  $m \geq \lceil \frac{k}{2} \rceil$  and  $n \geq \lceil \frac{k-1}{2} \rceil$ .*

In that paper the case when  $\lambda$  and at least one of  $m$  and  $n$  are odd was also studied. In particular Truszczyński showed (using a simple counting argument) that if  $\lambda$  is odd,  $m \geq n$ , and either:  $m$  and  $k$  are even,  $n$  is odd and  $k-1 > \lambda n$ ;  $m$  is odd,  $n$  and  $k$  are even and  $k-1 > \lambda m$ ; or  $m$  and  $n$  are odd, and  $k-1 > \lambda n$ , then there is no  $(K_{m,n}^\lambda, P_k)$ -design. This discovery led him to conjecture that except in these cases there exists a  $(K_{m,n}^\lambda, P_k)$ -design if and only if  $\lambda mn \equiv 0 \pmod{k-1}$ ,  $m \geq \lceil \frac{k}{2} \rceil$  and  $n \geq \lceil \frac{k-1}{2} \rceil$ .

Ushio [69] and Ushio and Tsuruno [70] studied resolvable  $(K_{a_1, a_2, \dots, a_m}, P_k)$ -designs.

**Theorem 2.14.** (Ushio) *There exists a resolvable  $(K_{m,n}, P_3)$ -design if and only if  $m+n \equiv 0 \pmod{3}$ ,  $m \leq 2n \leq 4m$  and  $3mn \equiv 0 \pmod{2(m+n)}$ .*

**Theorem 2.15.** (Ushio and Tsuruno) *There exists a resolvable  $(K_{(m,n)}^\lambda, P_3)$ -design if and only if  $mn \equiv 0 \pmod{3}$  and  $\lambda(m-1)n \equiv 0 \pmod{4}$ .*

Some time later Yu [72] substantially extend this work.

**Theorem 2.16.** (Yu) *For  $k \geq 4$ , there exists a  $(K_{(m,n)}^\lambda, P_k)$ -design if  $mn \equiv 0 \pmod{k}$ ,  $\lambda(m-1)kn \equiv 0 \pmod{2(k-1)}$  and either  $n \equiv 0 \pmod{k}$  or  $m \equiv 0 \pmod{k}$ .*

**Corollary 2.17.** *If  $k$  is a prime, there exists a  $(K_{(m,n)}^\lambda, P_k)$ -design if and only if  $mn \equiv 0 \pmod{k}$  and  $\lambda(m-1)kn \equiv 0 \pmod{2(k-1)}$ .*

Using the above results Yu is also able to show that the necessary conditions are sufficient when  $m = 2, 3$ .

**Corollary 2.18.** (Yu) (a) *There exists a  $(K_{n,n}^\lambda, P_k)$ -design if and only if  $2n \equiv 0 \pmod{k}$  and  $\lambda kn \equiv 0 \pmod{2(k-1)}$ .*

(b) *There is a  $(K_{n,n,n}^\lambda, P_k)$ -design if and only if  $3n \equiv 0 \pmod{k}$  and  $\lambda n \equiv 0 \pmod{k-1}$ .*

Caro and Schönheim [14] gave necessary and sufficient conditions for the existence of a  $(T, P_k)$ -design, where  $T$  is a tree.

We now turn our attention to the existence of  $(K_n^\lambda, F)$ -designs where  $F$  is a linear forest other than a path or the disjoint spanning union of isomorphic

paths and begin by considering  $(G, F)$ -designs. Bialostocki and Roditty [11] gave necessary and sufficient conditions for a  $(G, 3P_2)$ -design showing that with a finite number of exceptions, the design exists if and only if the number of edges in  $G$  is divisible by 3 and  $\Delta(G) \leq \frac{|E(G)|}{3}$ . This result was improved substantially by Alon [2].

**Theorem 2.19.** (Alon) *There exists a  $(G, tP_2)$ -design if and only if  $|E(G)| \equiv 0 \pmod{t}$  and  $t\chi'(G) \leq |E(G)|$ .*

*Proof.* The necessity is obvious. For the sufficiency, first observe that if in some graph there are two disjoint matchings of sizes  $m$  and  $n$ , where  $m > n$ , then there exist disjoint matchings of sizes  $m - 1$  and  $n + 1$ . Since,  $G$  has an edge-partitioning into  $\chi'(G)$  matchings, applying the observation to pairs of matchings yields the result.

**Corollary 2.20.** *There exists a  $(K_n^\lambda, tP_2)$ -design if and only if  $\lambda n(n - 1) \equiv 0 \pmod{2t}$  and  $t \leq \lfloor \frac{n}{2} \rfloor$ .*

For the case  $\lambda = 1$  this result was proven much earlier by Folkman and Fulkerson [23]. The design is *cyclic* if there is an  $n$ -cycle  $\sigma$  which permutes the vertices and also maps the set of all matchings onto itself. Recently Rees [56] has shown that in fact it is always possible to find cyclic  $(K_n^\lambda, tP_2)$ -designs when  $t < \frac{n}{2}$ . The case when  $t = \frac{n}{2}$  ( $n$  even) was dealt with earlier by Hartman and Rosa [27] who showed the designs exist except when  $n$  is at least 8 and a power of 2, in which case they do not exist.

Alon's result also settles the question for complete bipartite graphs.

**Corollary 2.21.** *There exists a  $(K_{n,m}^\lambda, tP_2)$ -design if and only if  $\lambda nm \equiv 0 \pmod{t}$  and  $t \leq \min\{m, n\}$ .*

Favaron, Lonc and Truszczyński [21] considered other subgraphs; in particular they considered the linear forest  $P_2 \cup P_3$ .

**Theorem 2.22.** (Favaron, Lonc and Truszczyński) *There exists a  $(G, P_2 \cup P_3)$ -design if and only if  $|E(G)| \equiv 0 \pmod{3}$ ,  $\Delta(G) \leq \frac{2}{3}|E(G)|$ ,  $c(G) \leq \frac{|E(G)|}{3}$  (where  $c(G)$  is the number of components of  $G$  with an odd number of edges), the edges of  $G$  cannot be covered by two adjacent vertices and  $G$  is not one of six specified graphs.*

**Corollary 2.23.** *There exists a  $(K_n^\lambda, P_2 \cup P_3)$ -design if and only if  $n \geq 5$  and  $\lambda n(n - 1) \equiv 0 \pmod{6}$ .*

Kotzig [45] claimed that if  $G$  is a 3-regular graph, there exists a  $(G, P_4)$ -design if and only if  $G$  has a 1-factor. (To see this observe that the number of paths of length 3 is the size of a 1-factor.) He asked for necessary and sufficient



conditions for a  $(G, P_{2n+2})$ -design when  $G$  is  $(2n+1)$ -regular. The only result we have along these lines is that of Jacobson, Truszczyński and Tuza [42]. (We do note, however, that Jünger, Reinelt and Pulleyblank [43] have asked if there exists a  $(G, P_4)$ -design whenever  $G$  is a simple planar 2-edge connected bipartite graph satisfying  $4|E(G)| \equiv 0 \pmod{3}$ .)

**Theorem 2.24.** (Jacobson, Truszczyński and Tuza) *There exists a  $(G, P_5)$ -design for every 4-regular bipartite graph  $G$ .*

Closely related to this is the work of Bondy [12] on perfect path double covers of graphs. He conjectured that for every graph  $G$  there is a decomposition of  $G^2$  into paths in which every vertex is the end-vertex of exactly two paths (paths of length zero are assumed to have two identical end-vertices). This conjecture was proven by Li [48]. Bondy further defined a  $k$ -regular path double cover of a graph  $G$  which in our terminology is a  $(G^2, P_k)$ -design, calling it *perfect* if every vertex is the end-vertex of exactly two paths. He asked for which  $G$  and  $k$  such designs exist.

**Theorem 2.25.** (Bondy) *If  $G$  is a connected simple graph, there exists a  $(G^2, P_3)$ -design if and only if  $G$  is not a tree with a 1-factor.*

**Theorem 2.26.** (Bondy) *If  $G$  is a connected simple graph, there exists a perfect  $(G^2, P_3)$ -design if and only if  $G$  is unicyclic and each vertex of degree 2 lies on the cycle.*

Bondy extended Kotzig's conjecture by conjecturing that if  $G$  is a simple  $k$ -regular graph, then there exists a perfect  $(G^2, P_{k+1})$ -design. In support of this conjecture he proved its validity in the case  $k = 3$ .

Turning now to other linear forests, Ruiz [63] has shown that there is a  $(K_{2n}, F)$ -design for any linear forest  $F$  with  $n$  edges. His result is proven by choosing the "right" set of edges from the path  $F^1$  (as described in Walecki's construction given in Theorem 2.1) to yield a subgraph isomorphic to  $F$  in which the differences on each edge are distinct, and then using it to generate the design. The work of Ruiz was generalized by Martinova [52] who used the same idea to show that there is a  $(K_{n,n}, F)$ -design where  $F$  is any linear forest on  $n$  edges.

### 3. $(G, \mathcal{F})$ -designs, where each $F \in \mathcal{F}$ is a linear forest.

When Tarsi [65] obtained his results on  $(K_n^\lambda, P_k)$ -designs, he was easily able to extend the work to address the following more general question:  
*Given a collection of paths with the property that each has at most  $n$  vertices*

and the sum of the lengths is  $\binom{n}{2}$ , is there a decomposition of  $K_n$  into precisely these paths?

This question (in the case  $\lambda = 1$ ) was also posed by Slater [64] who called a graph  $G$  *path-arboreal* if for a multiset  $\{a_i : 1 \leq i \leq r\}$  of  $r$  positive integers satisfying  $2 \leq a_i \leq |V(G)|$  and  $\sum_{i=1}^r (a_i - 1) = |E(G)|$ , there exists a  $(G, \mathcal{F})$ -design, where  $\mathcal{F} = \{P_{a_i} : 1 \leq i \leq r\}$ .

**Theorem 3.1.** (Tarsi) *If  $(n - 1)\lambda$  is even, and  $\{a_i : 1 \leq i \leq r\}$  is a multiset of  $r$  positive integers satisfying  $2 \leq a_i \leq n - 2$  and  $\sum_{i=1}^r (a_i - 1) = \lambda \binom{n}{2}$ , then there exists a  $(K_n^\lambda, \mathcal{F})$ -design, where  $\mathcal{F} = \{P_{a_i} : 1 \leq i \leq r\}$ .*

Note that this theorem not only imposes conditions on the parity of  $\lambda$  and  $n$  but restricts the lengths of the paths. Improvements were obtained by Ng [54], who considered only the case  $\lambda = 1$  and  $n$  odd. His main result allows all but Hamilton paths, although he did under certain conditions also construct designs in which some of the paths were Hamiltonian.

**Theorem 3.2.** (Ng) *If  $n$  is odd and  $\{a_i : 1 \leq i \leq r\}$  is a multiset of  $r$  positive integers satisfying  $2 \leq a_i \leq n - 1$  and  $\sum_{i=1}^r (a_i - 1) = \binom{n}{2}$ , then there exists a  $(K_n, \mathcal{F})$ -design, where  $\mathcal{F} = \{P_{a_i} : 1 \leq i \leq r\}$ .*

Truszczyński [67] investigated the question: *Is  $K_{m,n}^\lambda$  path-arboreal?*

**Theorem 3.3.** (Truszczyński) *If  $\{a_i : 1 \leq i \leq r\}$  is a multiset of  $r$  positive integers satisfying  $a_i \geq 2$ ,  $\sum_{i=1}^r (a_i - 1) = \lambda mn$ , and either*

1.  $m > n$  and  $a_i \leq 2n, i = 1, 2, \dots, r$ ,
2.  $m = n = 2$  and  $a_i \leq 4, i = 1, 2, \dots, r$ , or
3.  $m = n \geq 2$  and  $a_i \leq 2n - 4, i = 1, 2, \dots, r$ , then there exists a  $(K_{m,n}^\lambda, \mathcal{F})$ -design, where  $\mathcal{F} = \{P_{a_i} : 1 \leq i \leq r\}$ .

Fink and Straight [22] define a graph  $G$  to be *path-perfect* if there exists a  $(G, \mathcal{F})$ -design, where  $|E(G)| = \binom{n}{2}$  and  $\mathcal{F} = \{P_2, P_3, \dots, P_n\}$ . The first result they state is not difficult to verify.

**Theorem 3.4.** *The complete graph  $K_n, n \geq 2$ , is path-perfect.*

*Proof.* When  $n$  is odd, there is a  $(K_n, C_n)$ -design (Lucas [51]) and  $C_n \cong P_{k+1} \cup P_{n-k+1}, 1 \leq k \leq \frac{n-1}{2}$ . When  $n$  is even, there is a  $(K_n, P_n)$ -design (Theorem 2.1) and  $P_n \cong P_{k+1} \cup P_{n-k}, 1 \leq k \leq \frac{n}{2} - 1$ .

In that paper they also claimed that both  $K_{r,2r-1}$  and  $K_{r,2r+1}$  are path-perfect, and considered the existence of regular path-perfect graphs. Considerable progress on the question of path-perfect complete bipartite graphs has been made independently by Truszczyński [67] and Chilakamurri and Hamburger [16], although only in the case when  $s$  is even.

**Theorem 3.5.** (Chilakamurri and Hamburger; Truszczyński) *If  $s > t$  and  $s$  is even,  $K_{s,t}$  is path-perfect if and only if  $2st = n(n+1)$  for some integer  $n$  and  $2t \geq n$ .*

More generally, Zaks and Liu [74] obtained results on decomposing complete bipartite graphs into  $r$  paths, one of each even length from 2 to  $2r$ , and one of each odd length from 1 to  $2r - 1$ .

**Theorem 3.6.** (Zaks and Liu) *There exists a  $(K_{n,n}, \{P_2, P_4, \dots, P_{2n}\})$ -design for all  $n$ .*

**Theorem 3.7.** (Zaks and Liu) *There exists a  $(K_{n,n+1}, \{P_3, P_5, \dots, P_{2n+1}\})$ -design for all odd  $n$ .*

This latter result was extended by Chilakamurri [15] who showed that these designs also exist for all even  $n$  except  $n = 2$ , in which case the design does not exist.

**Theorem 3.8.** (Chilakamurri) *There exists a  $(K_{n,n+1}, \{P_3, P_5, \dots, P_{2n+1}\})$ -design for all even  $n$  except  $n = 2$ .*

By being less restrictive on the linear forests in  $\mathcal{F}$  (that is, by saying only that the subgraphs in the decomposition must be members of a set  $\mathcal{F}$  of linear forest) Lonc [50], Favaron, Lonc and Truszczyński [21] and Truszczyński [68] have obtained several results.

**Theorem 3.9.** (Lonc) *If  $\gcd(k, t) = 1$ , there exists an integer  $c = c(k, t)$  so that for every graph  $G$  satisfying  $\delta(G) \geq c$ , there is a decomposition of  $G$  into paths, each of length  $k$  or  $t$ .*

**Theorem 3.10.** (Favaron, Lonc and Truszczyński) *If  $G$  is a connected graph and  $G \not\cong K_3$ , then there is a decomposition of  $G$  into paths of lengths 2 and 3 if and only if  $G$  is not a tree with all vertices of odd degree.*

**Theorem 3.11.** (Truszczyński) *Every sufficiently large graph  $G$  satisfying  $|E(G)| \equiv 0 \pmod{k}$  and  $\Delta(G) \leq \frac{2}{k}|E(G)|$  has a decomposition into linear forests each containing precisely  $k$  edges.*

Note that these results do not specify the number of linear forests of each type. The only substantial result we have in which there are at least two types and the number of each is specified is due to Bermond, Heinrich and Yu [7] and concerns resolvable designs.

**Theorem 3.12.** (Bermond, Heinrich and Yu) *There is a resolvable decomposition of  $K_n^\lambda$  into  $s$  forests each isomorphic to  $\frac{n}{2}P_2$  and  $t$  forests each isomorphic*

to  ${}^n P_k$  when  $st \neq 0$ , if and only if  $n \equiv 0 \pmod{2}$ ,  $n \equiv 0 \pmod{k}$  and  $ks + 2t(k-1) = \lambda k(n-1)$ .

#### 4. Related questions.

In this section we very briefly mention an assortment of somewhat related questions. We begin with directed path-designs.

The only directed  $(G, F)$ -designs we will describe are those in which  $G$  is a symmetric directed graph (that is,  $G$  contains the arc  $(a, b)$  exactly when it contains the arc  $(b, a)$ ). The first result is an immediate corollary to Theorem 2.1.

**Theorem 4.1.** *There exists a  $(DK_n, \vec{P}_n)$ -design when  $n$  is even.*

The case  $n$  odd is considerably more difficult and it was not until 1980 that Tillson [66] constructed examples of such designs. (The question had been raised by Bermond and Faber [5] who attributed it to E.G. Strauss.)

**Theorem 4.2.** *(Tillson) There exists a  $(DK_n, \vec{P}_n)$ -design when  $n$  is odd and  $n \neq 3, 5$ , in which cases the designs do not exist.*

The existence of  $(DK_n, \vec{P}_k)$ -designs has not otherwise been considered except in the case of resolvable  $(DK_n, \vec{P}_k)$ -designs. Yu [73] has completely settled the question of their existence provided  $k \neq 3, 5$ . (And it is no accident that we have the same constraint on path length as in the previous theorem.)

**Theorem 4.3.** *(Yu) There exists a resolvable  $(DK_n^\lambda, \vec{P}_k)$ -design,  $k \neq 3, 5$ , exactly when  $n \equiv 0 \pmod{k}$  and  $n \equiv 1 \pmod{k-1}$ .*

Truszczyński [67] also studied directed designs.

**Theorem 4.4.** *(Truszczyński) If  $m \geq n$ , there exist a  $(K_{m,n}^\lambda, \vec{P}_k)$ -design if and only if  $2\lambda mn \equiv 0 \pmod{k-1}$ ,  $m \geq \lceil \frac{k}{2} \rceil$  and  $n \geq \lceil \frac{k-1}{2} \rceil$ .*

Heinrich and Nonay [29] considered the question of the existence of  $(K_n^2, P_n)$ -designs with the property that any two of the paths have exactly one edge in common. Designs for  $n \leq 20$  (but  $n \neq 4, 7$ , in which case the designs do not exist) were given and recursive constructions described. The construction was later improved by Horton and Nonay [35] who described a more general "product" technique. (We remark that such designs were used as a tool in partially answering a question of Hering [32] who asked for  $(K_n^2, C_{n-1})$ -designs with this intersection property.) An analogous question can be asked in the directed case where we ask for  $(DK_n, \vec{P}_k)$ -designs in which any two paths have exactly one oppositely directed arc in common (see [29]).

A  $(G, F)$ -design is *self-complementary* if  $F$  is a self-complementary graph (that is, if  $K_{|V(F)|}$  has a decomposition into exactly two copies of  $F$ ) and if on replacing each copy of  $F$  in the  $(G, F)$ -design by its complement, the result is another  $(G, F)$ -design. (For a survey of self-complementary graph decompositions see [59].) The only self-complementary linear forest is the path  $P_4$ . Granville, Moisiadis and Rees [24] constructed self-complementary  $(K_n, P_4)$ -designs, showing that they exist exactly when  $n \equiv 1 \pmod{3}$ .

Alspach [3] showed that the collection of all paths  $P_3$  with vertices from an  $n$ -set can be partitioned into copies of  $K_n$  if and only if  $n \equiv 0, 1 \pmod{4}$ . The analogous result for directed paths of length 2 was obtained by Heinrich [28].

Horton [34] constructed a  $(K_{2n}, P_{2n})$ -design  $\mathcal{A}$  with an associated  $(K_{2n}, nP_2)$ -design  $\mathcal{B}$  with the property that each  $nP_2$  (1-factor) in the decomposition  $\mathcal{B}$  has exactly one edge from each of the  $n$  paths in the decomposition  $\mathcal{A}$  for all values of  $n$  satisfying  $\gcd(n, 30) = 1$ .

Li [49] reports on progress made concerning the question of classifying all graphs with the property that their edges can be decomposed into paths each of length at least 3 (the question being attributed to W. R. Pulleyblank).

One variant of the "basic question" stated in Section 1 is obtained by either allowing some edges of  $G$  not to be used in the decomposition or by requiring that all edges be used but allowing some to be used more than once. More specifically, let  $G$  and  $F$  be graphs and define  $p(G, F)$  to be the maximum number of edge-disjoint copies of  $F$  in  $G$ , and  $c(G, F)$  to be the minimum number of copies of  $F$  required to contain every edge of  $G$  at least once. The first of these is called the *packing number* and the second the *covering number*. (There have been many papers on determining these numbers and the reader is referred to the (now somewhat old) survey articles by Beineke [4] and Harary [25].) The only results specifically on linear forests are due to Roditty [60, 61, 62] who determined the packing and coverings numbers in all cases when  $G \cong K_n$  and  $F$  is any linear forest on at most 5 vertices or  $F$  is the path of length 5. (He also obtained results for several other forests.) We note that using some of the results of Section 2 and the recursive construction technique it is not difficult to show that  $p(K_n, F) = \lfloor \frac{\binom{n}{2}}{|E(F)|} \rfloor$  and  $c(K_n, F) = \lceil \frac{\binom{n}{2}}{|E(F)|} \rceil$ , where  $F$  is any linear forest on at most 6 vertices.

M. L. Yu (personal communication) has recently asked the following question. Colour the edges of  $P_{2k+1}$  red and blue so that there are  $k$  edges of each colour, and colour the edges of  $K_n^2$  red and blue so that there is a red and a blue copy of  $K_n$ . If  $n(n-1) \equiv 0 \pmod{2k}$ , is there a decomposition of the coloured complete graph into copies of the coloured path?

Finally there has been much study of path numbers, path-covering numbers

and linear arboricity. The *linear arboricity* of a graph  $G$  is the minimum number of linear forests required to decompose the graph. Harary [25] conjectured that for any graph  $G$  the linear arboricity is between  $\lceil \frac{\delta(G)}{2} \rceil$  and  $\lceil \frac{\delta(G)+1}{2} \rceil$ . The *path number* of  $G$  is the minimum number of paths required to partition  $E(G)$ , and the *path-covering number* is the minimum number of paths necessary to cover the edges of  $G$ . In 1960, Gallai conjectured by Chung [17] that the path-covering number is also at most  $\lceil \frac{n}{2} \rceil$ . These questions are of quite a different flavour than those we have so far considered and for that reason we will leave the reader to consult appropriate source. (Similar numbers have also been investigated for directed graphs and tournaments.)

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